

Consistency Techniques over Real Numbers

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Reduction and Propagation

- Compute the solution set of $\{y = x^2, x \geq y + 1\}$ in the box $[-10, 10]^2$
- Propagation = sequence of reductions

$$y = x^2 \implies y \in [0, 10]$$

$$x \geq y + 1 \implies x \in [1, 10]$$

$$y \leq x - 1 \implies y \in [0, 9]$$

$$y = x^2 \implies y \in [1, 9]$$

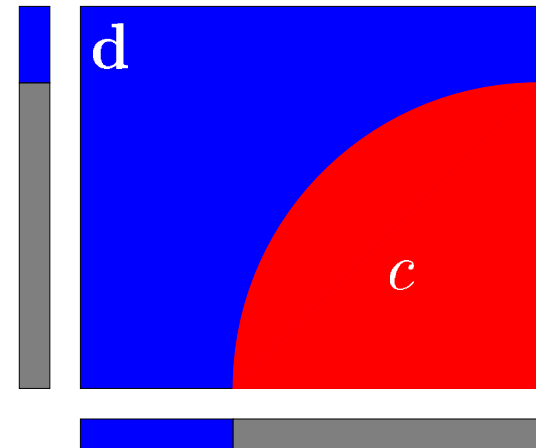
$$y = x^2 \implies x \in [1, 3]$$

...

Projections of a Constraint

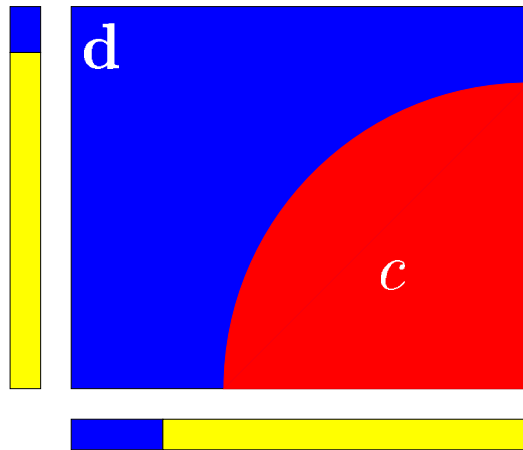
- Constraint $c(x_1, \dots, x_n) \subseteq \mathbb{R}^n$
- Box $d = d_1 \times \dots \times d_n$
- Projection of c over x_i

$$\{a_i \in d_i \mid \exists a_1 \in d_1, \dots, \\ \exists a_{i-1} \in d_{i-1}, \\ \exists a_{i+1} \in d_{i+1}, \dots, \\ \exists a_n \in d_n : \\ c(a_1, \dots, a_n)\}$$



Approximation of Projections

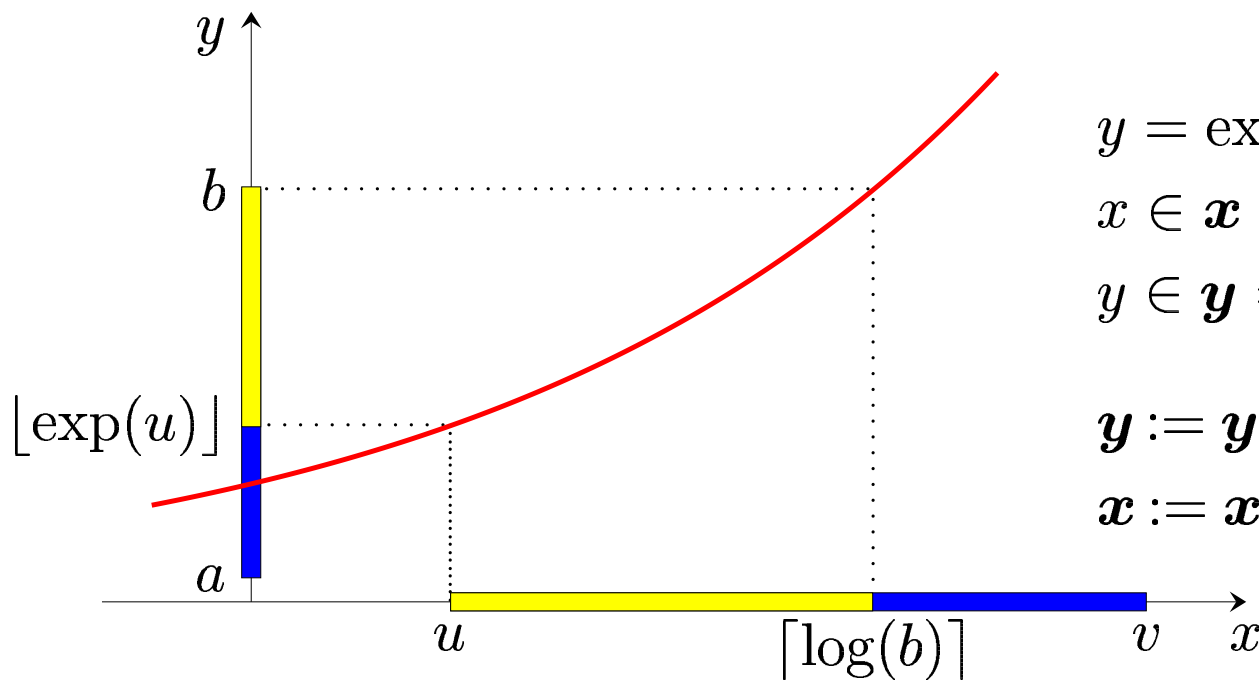
- Projections cannot be represented by floating-point intervals
- Approximation



- Reduction is computation of some yellow box

Reduction: Inversion

- Cleary 87
- Reduce the domains of x and y st. $y = \exp(x)$



$$y = \exp(x)$$

$$x \in \mathbf{x} = [u, v]$$

$$y \in \mathbf{y} = [a, b]$$

$$\mathbf{y} := \mathbf{y} \cap \exp([u, v])$$

$$\mathbf{x} := \mathbf{x} \cap \exp^{-1}([a, b])$$

Inversion for Complex Constraints

- Cleary 87
- Decomposition into primitives

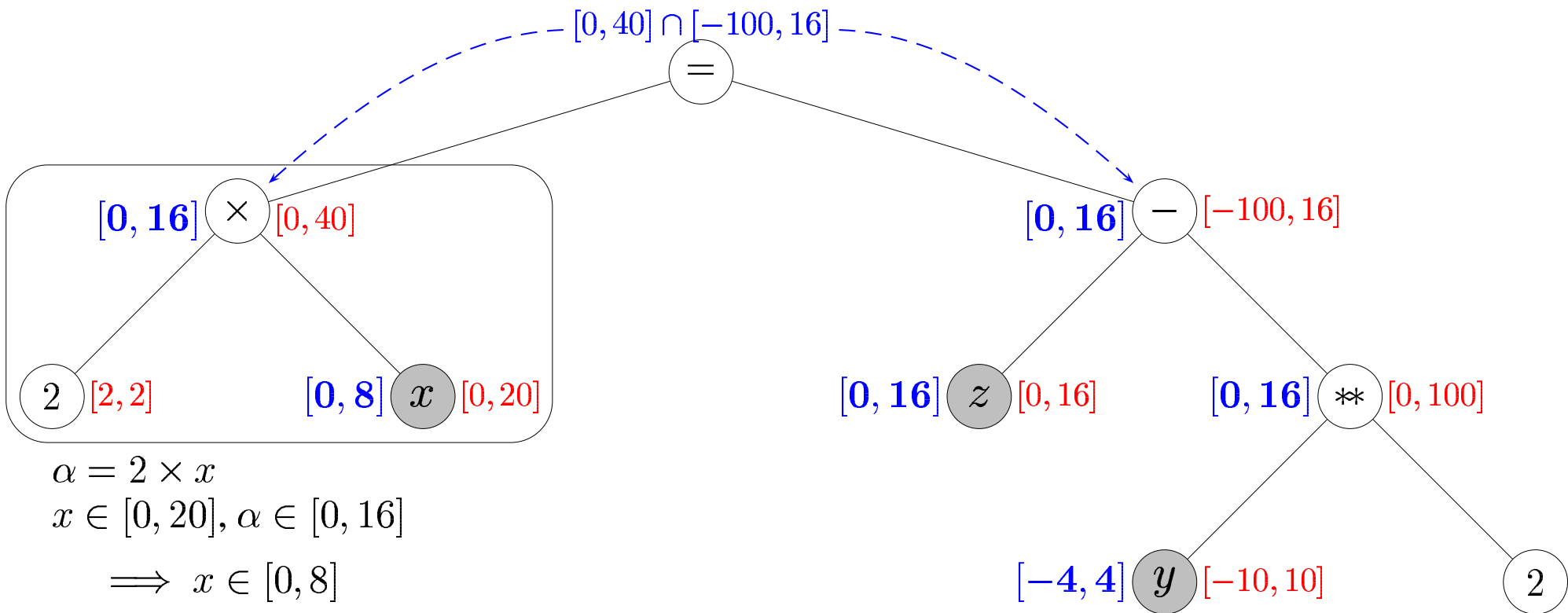
$$2x = z - y^2 \iff \begin{cases} \alpha = 2 \times x \\ \beta = y^2 \\ \alpha = z - \beta \end{cases}$$

for some $\alpha, \beta \in \mathbb{R}$

- Slow compilation and propagation processes

Inversion for Complex Constraints

- Benhamou *et al.* 99



Dependency Problem

- The inversion technique can be weak

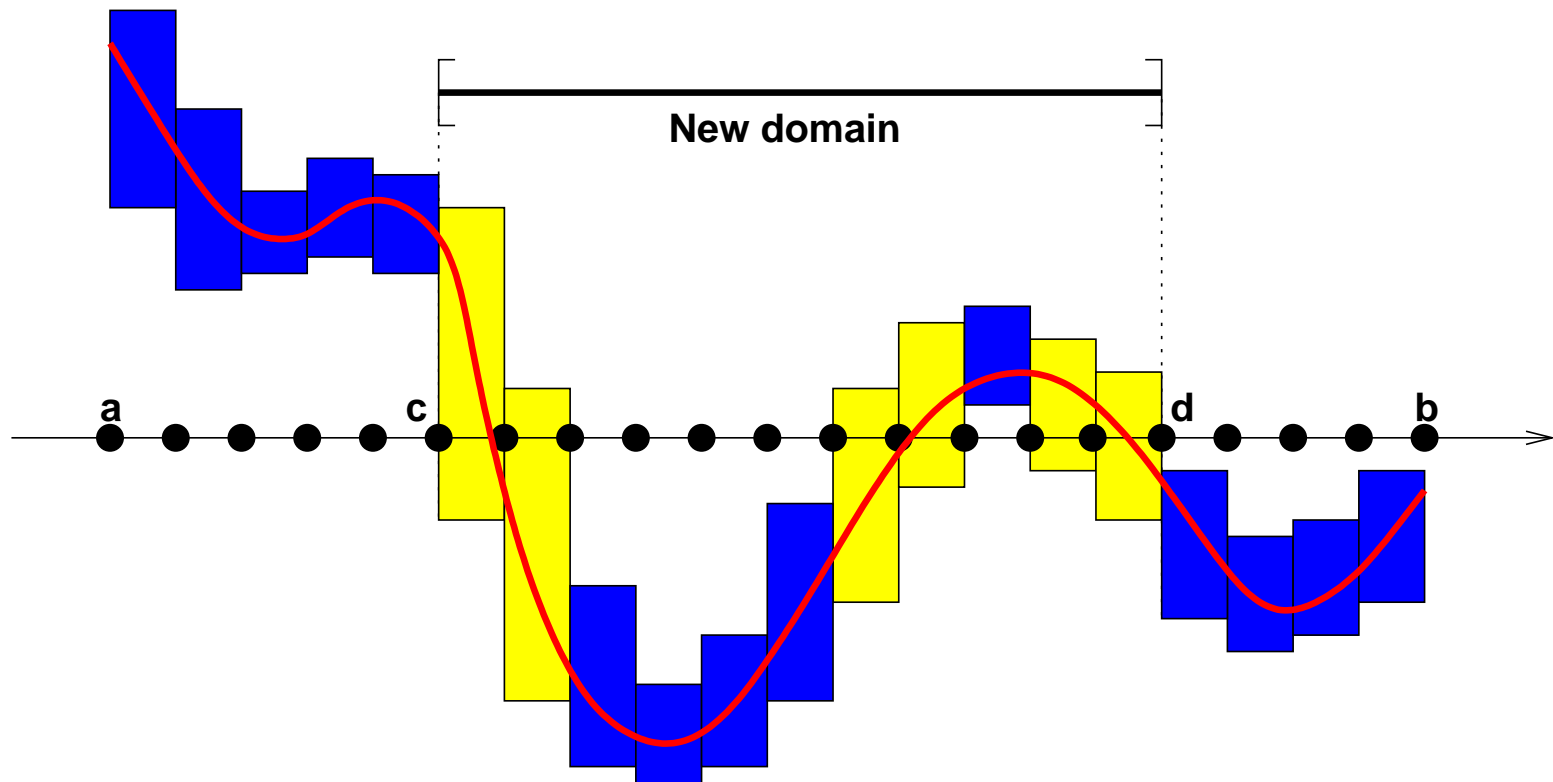
$$x + x = 0, x \in \mathbf{d} = [-1, 1]$$

$$\begin{aligned} \mathbf{d} &:= \mathbf{d} \cap (-\mathbf{d}) \\ &:= [-1, 1] \cap [-1, 1] \\ &:= [-1, 1] \end{aligned}$$

- Solutions
 - symbolic transformation
 - bisection-evaluation process

Reduction: Box-consistency

- Benhamou & McAllester & Van Hentenryck 94
- Reduce the domain of x st. $f(x) = 0$



Box-co. for Multivariate Constraints

- Consider a constraint $f(x_1, \dots, x_n) = 0$ and an interval form f of f
- Given a box \mathbf{d} box-consistency is computed for the set of interval functions

$$\left\{ \begin{array}{l} f_1(x_1, \mathbf{d}_2, \mathbf{d}_3, \dots, \mathbf{d}_n) \\ f_2(\mathbf{d}_1, x_2, \mathbf{d}_3, \dots, \mathbf{d}_n) \\ \quad \quad \quad \ddots \\ f_n(\mathbf{d}_1, \mathbf{d}_2, \dots, \mathbf{d}_{n-1}, x_n) \end{array} \right.$$

Reduction Model

- Let \mathbb{I} be the set of closed intervals
- A **reduction operator** is a function $\mathbb{I}^n \rightarrow \mathbb{I}^n$ which is for all $x, y \in \mathbb{I}$
 - monotonic: $x \subseteq y \implies \theta(x) \subseteq \theta(y)$
 - contracting: $\theta(x) \subseteq x$
- Consistency techniques are monotonic and contracting
- These techniques are also **verified!**

Propagation Model

- Narin'yani 1983, Benhamou 1996, Apt 1998
- Lemma: given a finite set of reduction operators and a domain d the **greatest common fixed-point** of the operators included in d exists
- Propagation = application of operators
 - fair strategy \implies **convergence**
 - finiteness of domain \implies **termination**
 - **independence** wrt. the strategy

Fixed-point Propagation Algorithm

$S := \{0_1, \dots, 0_n\}$

Box d

repeat

 choose 0 in S

 save := d

$d := O(d)$

 if $d = \text{save}$ then

$S := S - \{0\}$

 else

$S := \{0_j \mid 0_j \text{ depends on a modified domain}\}$

until S is empty

Solver Cooperation

1. Generation of reduction operators
 - Inversion
 - Box-consistency
 - Interval methods
 - Simplex
2. Propagation over the set of operators
 - good strategy (choose) \implies efficiency

A Strategy

`S := empty`

`for each constraint C do`

`if there exists X in C occurring once`

`then S := S U { inversion(c) }`

`for each variable X in C do`

`if X occurs more than once in C`

`then S := S U { box(C,X) }`

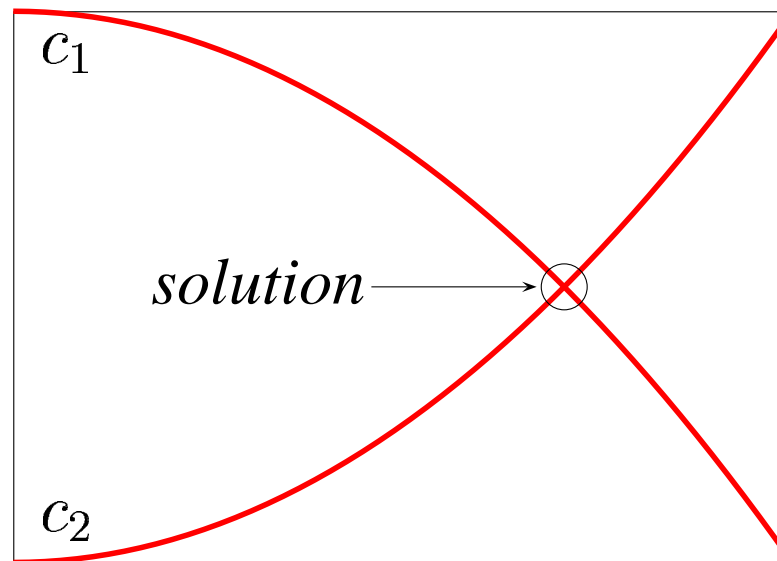
`od`

`if a square system of equations E can be computed`

`then S := S U { IntervalNewton(E) }`

Locality Problem

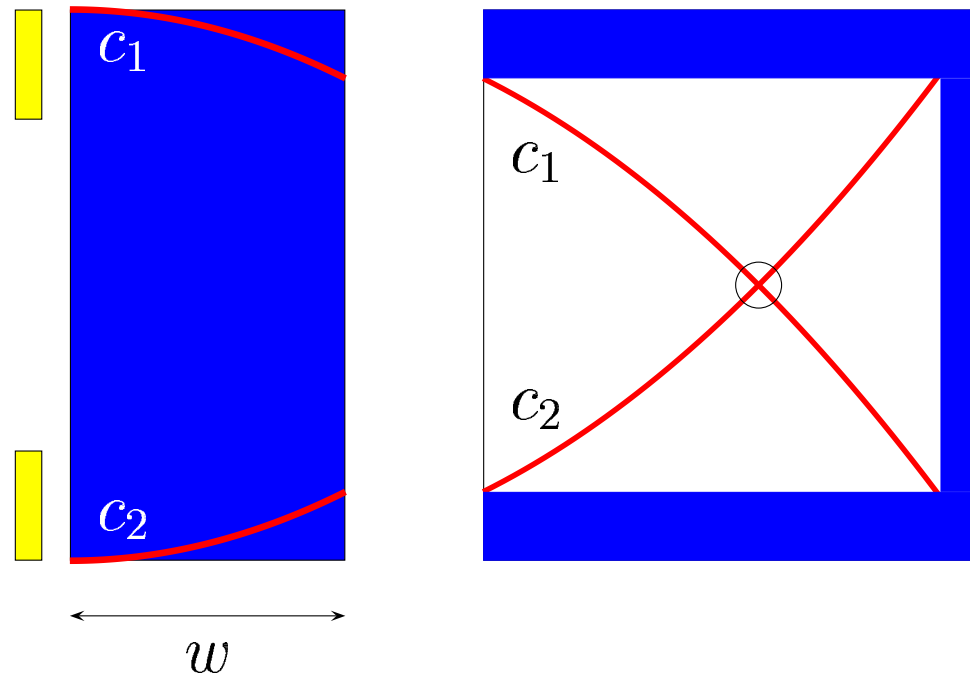
- Consider a difficult problem for inversion or box-consistency



- The intersection of projections is weaker than the projection of the intersection

Reduction: 3B-consistency

- Lhomme 1993
- Prove the inconsistency of a sub-domain at one bound of a variable domain



Bisection Algorithm

- Classical algorithm
 - precision of boxes: 10^{-8}
 - choice of variable: round-robin
 - bisection in 3 parts
- Reductions = cooperation of
 - I: inversion
 - B: box-consistency
 - N: interval Newton

Results

Problem	v	n	I-B-N	I-B	B-N	I-N
BB	10	1	0.1	0.2	0.1	3.4
	20	1	0.2	0.5	0.2	?
	40	1	0.5	1.0	0.5	?
	80	1	1.1	2.1	1.1	?
	160	1	2.2	4.1	2.2	?
MC 1	10	1	0.0	0.1	0.1	0.1
	20	1	0.1	0.4	0.7	0.4
	40	1	0.5	1.7	5.2	1.4
	80	1	2.4	6.9	41.4	6.4
Neuro	8	8	46.1	178.9	141.9	46.1
Kin 2	8	10	25.0	1175.2	53.3	19.2
Kin 1	12	16	0.7	110.9	2.0	0.7

Conclusion

- Consistency techniques are used in projects
 - automatic control (CNRS group on set-based methods)
 - computer-aided design (EPFL, RNTL CO2)
 - global optimization (COCONUT project)
 - image synthesis (PRIAMM project)
- Research directions: new types of constraints, specific propagation algorithms, inner approximations, quantified constraints, mixed-problems...